## PRANGYA PARIDA, University of Ottawa

Cover-free families on graphs

A family of subsets of [t] is called a *d*-cover-free family (*d*-CFF) if no subset is contained in the union of any *d* others. We denote by t(d, n) the minimum *t* for which there exists a *d*-CFF of [t] with *n* subsets. t(1, n) is determined using Sperner's Theorem. For  $d \ge 2$ , we rely on bounds for t(d, n). Using the probabilistic approach, Erdös, Frankl, and Füredi proved  $3.106 \log(n) < t(2, n) < 5.512 \log(n)$ . Porat and Rothschild provided a deterministic polynomial-time algorithm to construct *d*-CFFs achieving  $t = O(d^2 \log(n))$ . Some upper bounds of t(2, n) (in some cases exact bounds) for some small values of *n* were provided by Li, van Rees, and Wei.

We extend the definition of 2-CFF to include a graph(G), called  $\overline{G}$ -CFF, where the edges of G specify the pair of subsets whose union must not cover any other subset. We denote by t(G) as the minimum t for which there exists a  $\overline{G}$ -CFF. Thus,  $t(K_n) = t(2, n)$ . We will discuss some classical results on CFFs, along with constructions of  $\overline{G}$ -CFFs. We prove that for a graph G with n vertices,  $t(1, n) \leq t(G) \leq t(2, n)$  and for an infinite family of star graphs with n vertices,  $t(S_n) = t(1, n)$ . We also provide constructions for  $\overline{P_n}$ -CFF and  $\overline{C_n}$ -CFF using a mixed-radix Gray code. This yields an upper bound for  $t(P_n)$  and  $t(C_n)$  that is smaller than the lower bound of t(2, n) mentioned above.