
PAUL TAYLOR, University of Manchester, UK
Computable Real Analysis without Set Theory or Turing Machines

The many schools of computable or constructive analysis accept without question the received notion of set with structure. They rein in the wild behaviour of set-theoretic functions using the double bridle of topology and recursion theory, adding encodings of explicit numerical representations to the epsilons and deltas of metrical analysis. Fundamental conceptual results such as the Heine–Borel theorem can only be saved by set-theoretic tricks such as Turing tapes with infinitely many non-trivial symbols.

It doesn't have to be like that.

When studying computable continuous functions, we should never consider uncomputable or discontinuous ones, only to exclude them later. By the analogy between topology and computation, we concentrate on open subspaces. So we admit $+$, $-$, \times , \div , $<$, $>$, \neq , \wedge and \vee , but not \leq , \geq , $=$, \neg or \Rightarrow . Universal quantification captures the Heine–Borel theorem, being allowed over *compact* spaces. Dedekind completeness can also be presented in a natural logical style that is much simpler than the constructive notion of Cauchy sequence, and also more natural for both analysis and computation.

Since open subspaces are defined as continuous functions to the Sierpiński space, rather than as subsets, they enjoy a “de Morgan” duality with closed subspaces that is lost in intuitionistic set-, type- or topos theories. Dual to \forall compact spaces is \exists over “overt” spaces. Classically, all spaces are overt, whilst other constructive theories use explicit enumerations or distance functions instead. Arguments using \exists and overtness are both dramatically simpler and formally dual to familiar ideas about compactness.